Duplicat Request Cache (DRC) for NFSv4 over TCP Rick Macklem, University of Guelph

Chet's Good Ole DRC for NFS over UDP

- Typically a Global LRU Cache
- Since UDP retries agressively
- \rightarrow throw away Hits while ''in progress''
- \rightarrow cache replies for non-idempotent RPCs
- \rightarrow reply from cache for Hits with cached replies
- some concern w.r.t. False Hits (reused xid values)

DRC for NFSv4 over TCP

- only retries after reconnect
- \rightarrow never on same TCP connection
- •1 minute to hours later
- \rightarrow other clients could still be active
- \rightarrow global LRU doesn't work well

- what about LRU per TCP connection?
- time based cache invalidation?
- size based cache invalidation?
- Ideally, when the client has processed the reply
- \rightarrow I invalidate when client side TCP acks receipt of reply (at least reply is in client rcv Q)
- Wired down entries for Seqid# Ops

False Hits

- all RPCs same # (Compound)
- time cached much longer
- \rightarrow risk of False Hit greater
- **‡** Therefore:
 - I assume a hit ''in progress'' \rightarrow False Hit allow multiple requests for same <cache key>

• use checksum on first bytes of NFS XDR

rightarrow I didn't use client IP# in cache key (DHCP lease expiry \rightarrow different IP#) When request arrives, match all:

if

same xid

‡ not "in progress" not same socket same length of NFS XDR same checksum for first ≤ 100 bytes of XDR 1 hit with no Seqid# Op in it → reply with cached entry **During processing of Compound:**

if non-idempotent Op → set flag
if Op uses a Seqid#
 if same seqid# as referenced entry → Hit
 free this entry
 if cached entry "in progress"
 drop request
 else
 reply from cached entry
 else if next seqid# in order
 free referenced cache entry
 wire down this cache entry
 else if first seqid#
 wire down this cache entry

End of Compound Processing:

if wired down OR (non-idempotent AND below Floodlevel) save reply in cache entry timestamp it

note TCP seq#

else

free cache entry

Send reply

• Certain error replies aren't cached: NFS4ERR_GARBAGE, NFS4ERR_BADXDR, NFS4ERR_BADSEQID, NFS4ERR_RESOURCE, NFS4ERR_STALECLIENTID, NFS4ERR_OLDSTATEID, NFS4ERR_BADSTATEID, NFS4ERR_GRACE, NFS4ERR_NOGRACE, NFS4ERR_MOVED, NFS4ERR_STALESTATEID, NFS4ERR_SERVERFAULT, NFS4ERR_DELAY - Unless a Seqid# Op **Cache Invalidation Happens When:**

- $\bullet \, For \, Seqid\# \, Op \; RPCs \rightarrow next \; Seqid\# \, Op \; processed$
- Others \rightarrow client TCP acks receipt of reply
- **OR** \rightarrow large timeout (12 hours)

NFSv2 and 3

- Over UDP \rightarrow same old DRC
- Over TCP, Not the same as NFSv4
 - no check for different TCP socket
 - cache key includes RPC#

Although my current code doesn't do so, I think:

‡ should drop request and TCP connection

Packrats: A Work starting to Progress

purely experimental, to see if agressive client side caching will improve perf over network interconnects will large "bandwidth * delay"

packrat threads do agressive client side data caching onto local storage when the server issues a Delegation to the client (whole file copies to the client as soon as the delegation is issued)

delegations are working, but the packrat threads aren't yet